

Research Article

International Journal of Scientific Research and Reviews

An Alternate C++ Programme for Total Dominator Chromatic Number of Paths and Cycles

J.Virgin Alangara Sheeba^{*} and A.Vijayalekshmi

Department of Mathematics, S.T.Hindu College, Nagercoi629002, Tamilnadu,India. Affiliated to Manonmaniam Sundaranar University, Abishekapatti, Thirunelveli-627012. TN, India

ABSTRACT

A total dominator coloring of a graph G=(V,E) without isolated vertices is a proper coloring together with each vertex in G properly dominates a color class. The total dominator chromatic number of G is a minimum number of color classes with additional condition that each vertex in G properly domi- nates a color class and is denoted by χ_{td} (G). In this paper we introduce an alternate C++ programme to find the total dominator chromatic number of paths and cycles.

MATHEMATICS SUBJECT CLASSIFICATION: 05C15, 05C69

KEYWORDS: Total dominator chromatic number, path graph, cycle graph.

*Corresponding author J.Virgin Alangara Sheeba

*Research Scholar, Department of Mathematics, S.T.Hindu College, Nagercoi629002, Tamilnadu, India. Affiliated to Manonmaniam Sundaranar University, Abishekapatti, Thirunelveli-627012.Tamil Nadu, India Email:vijimath.a@gmail.com

INTRODUCTION

In this paper we only consider paths and cycles. Further details in graph the- ory can be found in F.Harrary¹.

Let G=(V,E) be a graph with minimum degree at least one. For two vertices v_0 and v_n of a graph G,a v_0 - v_n walk is an alternate sequence of vertices and edges $v_0, e_1, v_1, \dots, v_n$ e_n , v_n such that the consecutive vertices and edges are incident. A path is a walk in which no vertex is repeated. The vertices v_0 and v_n are called the initial and terminal vertex respectively. A v_0 - v_n walk is closed if $v_0 = v_n$. A closed walk with no repeated vertices except the initial and terminal vertices is called a cycle. A path and cycle with n vertices are denoted by P_n and C_n respectively. A proper coloring of G is an assignment of colors to the vertices of G, such that adjacent vertices have different colors. The smallest number of colors for which there exists a proper coloring of G is called a chromatic number of G, and is denoted by $\chi(G)$. A total dominator coloring (td-coloring) of G is a proper coloring of G with extra property that every vertex in G properly dominates a color class. The total dominator chro- matic number is denoted by $\chi_{td}(G)$ and is defined by the minimum number of colors needed in a total dominator coloring of G. This concept was introduced by A.Vijayalekshmi in [2]. This notion is also referred as a smarandachely k-dominator coloring of G, $(k \ge 1)$ and was introduced by A.Vijayalekshmi in [2]. For an integer $k \ge 1$, a smarandachely k-dominator coloring of G is a proper coloring of G, such that every vertex in a graph G properly dominates a k color class. The smallest number of colors for which there exists a smarandachely k-dominator coloring of G is called the smarandachely kdominator chromatic number of G and is denoted by χ^s (G).

In a proper coloring C of a graph G, a color class of C is a set consisting of all those vertices assigned the same color. Let C be a minimum td-coloring of G. We say that a color class is called a non-dominated color class (n-d color lass) if it is not dominated by any vertex of G and these color classes are also called repeated color classes.

The total dominator chromatic number of paths and cycles were found in [3]. We have the following observation from [3]

Theorem A [3]

Let G be P_n or C_n . Then

$$\chi_{td}(P_n) = \chi_{td}(C_n) = \begin{cases} 2\left\lfloor \frac{n}{4} \right\rfloor + 2 & \text{if } n \equiv 0 \pmod{4} \\ 2\left\lfloor \frac{n}{4} \right\rfloor + 3 & \text{if } n \equiv 1 \pmod{4} \\ 2\left\lfloor \frac{n+2}{4} \right\rfloor + 2 & \text{Otherwise} \end{cases}$$

MAIN RESULT

We have to find the total dominator chromatic number of paths and cycles by using C++ programme by removing the corresponding rows and columns of the repeated colors from the adjacency matrix. The C++ programme is successfully compiled and run on C++ platform. The runtime test is included.

PROGRAM AS FOLLOWS

```
#include "stdafx.h"
#include <Windows.h>
#include <conio.h>
#include <iostream>
  using namespace std; int
  opt = 1;
int main()
    {
cout \ll "\n";
\operatorname{cout} << \operatorname{"ln"} <
Exit" << "\n";
cin >> opt;
  while (opt \leq 3\&\& opt > 0)
    {
  if (opt == 1 \parallel opt == 2) //---- if option 1 or 2
    {
```

```
int inpt;
if (opt == 1)
{
cout << "Enter the Value of Pn" << endl;
cin >> inpt;
}
else
{
cout << "Enter the Value of Cn" << endl;
cin >> inpt;
}
while (inpt \geq 11)
{
// dimensions
int N=inpt; // matrix row
int M=inpt; // matrix column
// dynamic allocation
int** a=new int*[N]; //logic matrix int** b=new int*[N]; //logic
matrix
int** mat = new int*[N]; //adjacency matrix
int** cc = new int*[N]; //adjacency matrix without repeating colors int** bb = new int*[N];//logic matrix
for (int i=0; i<N; ++i)
{
a[i] = new int[M]; b[i] = new int[M];
mat[i] = new int[M]; cc[i] = new int[M];
bb[i] = new int[M];
}
// variables int i, j, k; int n;
```

```
int g, h; int ii =0; int jj =0; int
d=0;
n=inpt;
HANDLE p=GetStdHandle(STD_OUTPUT_HANDLE);
if (n % 4== 1)
{
for (i =0; i<n; i++)
{
g=0;
h=3;
for (j =0; j<n; j++)
{
a[i][j] =j;
if (j == g \&\& g <= n-1)
{
b[i][j] = 0; g = j+4; bb[j][i] = 0;
}
else
if (j == h \& \& h \le n - 3 \& \& n - 1 \ge h)
{
b[i][j] =0; h=j+4; bb[j][i] =0;
}
else
{
b[i][j] =j;
bb[j][i] =j;
}
}
}
```

```
}
else
{
for (i =0; i<n; i++)
{
g=0;
h=3;
for (j =0; j<n; j++)
{
a[i][j] =j;
if (j == g \&\& g \le n - 3 \&\& n - 1! = g)
{
b[i][j] =0;
g=j+4; bb[j][i] =0;
}
else
if (j == h \& \& h \le n-1)
{
b[i][j] =0;
h=j+4;
bb[j][i] =0;
}
else
{
b[i][j] =j;
bb[j][i] =j;
}
}
}
}
```

```
// ----- END FORMING 2 ROW MATRIXES ------ for (i = 0; i < 1;
i++)
{
for (j =0; j < n; j++)
{
if (b[i][j] == 0)
{
SetConsoleTextAttribute(p, FOREGROUND_RED | FOREGROUND_INTENSITY);
//cout << b[i][j] << " ";</pre>
}
else
{
SetConsoleTextAttribute(p, FOREGROUND_INTENSITY | FOREGROUND_INTENSITY);
//cout << b[i][j] +1 << "";</pre>
if (b[i][j-1] == 0 \&\& b[i][j+1] != 0)
{
d = d + 1;
}
}
}
//cout << endl;</pre>
SetConsoleTextAttribute(p, FOREGROUND_INTENSITY | FOREGROUND_INTENSITY);
}
\operatorname{cout} \ll "\backslash n";
if (opt == 1)
{
```

```
"\n"; // GETTING VALUE FROM USER
}
else if (opt == 2)
{
"\n"; // GETTING VALUE FROM USER
}
// ----- LOGIC TO FORM MATRIX------ if (opt == 1)
{
for (i =0; i<n; i++)
{
for (j =0; j<n; j++)
{
if (b[i][j] == 0 || bb[i][j] == 0)
{
if (a[i][j] == i+1 | a[i][j] == i-1)
{
mat[i][j] =1;
SetConsoleTextAttribute(p, FOREGROUND_RED | FOREGROUND_INTENSITY);
cout << mat[i][j] << " ";
}
else
{
mat[i][j] =0;
SetConsoleTextAttribute(p, FOREGROUND_RED | FOREGROUND_INTENSITY);
cout << mat[i][j] << "";
}
}
else
```

```
{
if (a[i][j] == i+1 | a[i][j] == i-1)
{
mat[i][j] =1;
SetConsoleTextAttribute(p, FOREGROUND_INTENSITY | FOREGROUND_INTENSITY);
cout << mat[i][j] << " ";
}
else
{
mat[i][j] = 0;
SetConsoleTextAttribute(p, FOREGROUND_INTENSITY | FOREGROUND_INTENSITY);
cout << mat[i][j] << " ";
}
}
}
cout \ll "\n";
}
}
else if (opt == 2)
{
for (i =0; i<n; i++)
{
for (j = 0; j < n; j + +)
{
if (b[i][j] == 0 || bb[i][j] == 0)
{
if (a[i][j] == i+1 | a[i][j] == i-1 | a[i][j] == i+(n-1) |
    a[i][j] == i - (n - 1))
{
```

```
mat[i][j] = 1;
  SetConsoleTextAttribute(p, FOREGROUND_RED | FOREGROUND_INTENSITY);
cout << mat[i][j] << "";
}
else
{
mat[i][j] =0;
SetConsoleTextAttribute(p, FOREGROUND_RED | FOREGROUND_INTENSITY);
cout << mat[i][j] << "";
}
}
else
{
if (a[i][i] == i+1 |a[i][i] == i-1 |a[i][i] == i+(n-1) |
    a[i][j] == i - (n - 1))
{
mat[i][j] =1;
SetConsoleTextAttribute(p, FOREGROUND_INTENSITY | FOREGROUND_INTENSITY);
cout << mat[i][j] << "";
}
else
{
mat[i][j] = 0;
SetConsoleTextAttribute(p, FOREGROUND_INTENSITY | FOREGROUND_INTENSITY);
cout << mat[i][j] << " ";
}
}
}
cout \ll "\n";
}
```

}

 $cout \ll "\n";$

// -----END LOGIC TO FORM MATRIX------

SetConsoleTextAttribute(p, FOREGROUND_INTENSITY | FOREGROUND_INTENSITY);

// Adjecancy matrix without repeating colors

// Define $S = \{ Vertices of Pn and Cn receiving different colors \}$ and

// V-S = {Vertices of Pn and Cn receiving repeated colors }

// Assign different colors to the vertices in S. If vi and vj belongs to V-S

// and vi,vj are adjacent then we assign repeated colors 1,2 (say) to vi and vj $\,$

// respectively. Also vk belongs to V-S and vk is not adjacent to any vi belongs

// to V-S then assign either color 1 or color 2 to vk. If vi belongs to V-S then

 ${\ensuremath{\textit{//}}}$ remove the corresponding rows and columns of vi from the adjacency matrix.

cout <<"The Sub Matrix after removing the repeating colors are" << "\n"<< "\n";
for (i =0; i<n; i++)
{</pre>

```
for (j =0; j<n; j++)
{
    if (j == 0)
    {
        jj =0;
    }
    if (b[i][j] != 0 && bb[i][j] != 0)
    {
        cc[ii][jj] = mat[i][j];
        jj =jj + 1;
    }
    if (j == n-1 && bb[i][j] != 0)</pre>
```

```
{
ii =ii +1;
}
}
}
for (i =0; i<ii; i++)
{
for (j = 0; j < ii; j++)
{
cout << cc[i][j] << "";
}
cout \ll "\n";
} system("pause"); cout <<</pre>
"\n";
//cout << d;
cout \ll "\n";
if (n \% 4 == 1)
{
cout << "Number of (2X2) sub Matrices are "<< d-1<< " "<< "\n";
cout << "Number of (3X3) sub Matrices are "<< "1" << " '' << "\n" ;
}
else
{
cout << "Number of (2X2) sub Matrices are "<< d<< " "<< "\n";
cout << "Number of (3X3) sub Matrices are "<< "0" << ""<<"\mid n" << "\mid n";
}
if (n % 4 != 1)
{
```

```
for (i = 0; i < ii; i++)
{
for (j = 0; j < ii; j++)
{
if (cc[i][j] == 1 || j == i)
{
SetConsoleTextAttribute(p, FOREGROUND_RED | FOREGROUND_INTENSITY);
cout << cc[i][j] << "";
}
else
{
SetConsoleTextAttribute(p, FOREGROUND_INTENSITY | FOREGROUND_INTENSITY);
cout << cc[i][j] << "";
}
}
cout \ll "\n";
}
}
else
{
for (i = 0; i < ii; i++)
{
for (j = 0; j < ii; j++)
{
if (cc[i][j] == 1 || j == i || j = ii + 3 || i = ii - 3 \&\& j = ii - 3)
{
SetConsoleTextAttribute(p, FOREGROUND_RED | FOREGROUND_INTENSITY);
cout << cc[i][j] << "";
}
else
```

{

```
SetConsoleTextAttribute(p, FOREGROUND_INTENSITY | FOREGROUND_INTENSITY);
cout << cc[i][j] << "";
}
}
\operatorname{cout} \ll "\backslash n";
}
}
cout \ll "\n";
SetConsoleTextAttribute(p, FOREGROUND_INTENSITY | FOREGROUND_INTENSITY);
//system("pause");
cout << "Total Dominator Chromatic Number is" << "\n" << "2+[(2 *number of (2X2) matrices)+(3 *
 number of (3X3) matrices)]
\operatorname{cout} \ll "\backslash n";
if (n \% 4 == 1)
{
\operatorname{cout} \ll "\backslash n";
cout << "TOTAL DOMINATOR CHROMATIC NUMBER IS "<< 2+(2*d-2)+3<< "\n";
}
else
{
\operatorname{cout} \ll "\backslash n";
cout << "TOTAL DOMINATOR CHROMATIC NUMBER IS "<< 2+(2*d) << "\n";
}
system("pause");
// free ary and mat
for (int i=0; i<N; ++i)
delete[] mat[i];
```

delete[] mat;

```
delete[] cc[i]; delete[] cc; delete[]
  a[i]; delete[] a; delete[] bb[i];
 delete[] bb; delete[] b[i]; delete[] b;
  return main();
    }
                                                                                                                          \frac{1}{2} + \frac{1}
    {
 return 0;
    }
    }
cout << "please enter the above values "<< "\n";
  system("pause");
                                                                                                                                                                                             // restart program cout << "\n";</pre>
  return main();
 system("pause");
 return 0;
    }
```

RUNTIME TEST

```
please choose the below options
1. Path
2. Cycle
3. Exit
Enter the Value of Pn
11
The Adjacency Matrix for P11
with the repeating colors highlighted in red
                              55
                                 <mark>-</mark> 51 53
                5
                    5
      10
   5
   1
                                 S S
   SS
      SS
                   10
                              5
                5
                ĭ
   5
      5
                55
                   5
                                 1
0
                              0
1
The Sub Matrix after removing the repeating colors are
010000
100000
00100
00100
00001
00001
000010
Press any key to continue . . .
Number of (2x2) sub Matrices are 3
Number of (3x3) sub Matrices are Ø
      0
0
1
1
0
0
0
0
            2 2 2 2 2 2
                5555
   55555
2 2 2 3
Total Dominator Chromatic Number is
2 + [(2 * number of (2x2) matrices) + (3 * number of (3x3) matrices)]
TOTAL DOMINATOR CHROMATIC NUMBER IS 8
```

RUNTIME TEST

```
please choose the below options
    Path
Cycle
Exit
Enter the Value of Cn
The Adjacency Matrix for C13
with the repeating colors highlighted in red
                               5
                SS
                            5
                          5
   0
1
              30
     10
                         9 9
9 9
     S
                              5
   5
              Ø
1
                 1
                Ø
                555
                            1
Ø
1
                         5
   1
     2
              2
                               5
   5
              5
                          10
                               1
Ø
     5
     2
    Sub Matrix after removing the repeating colors are
The
     889488
        5515
           りうりのりょう
             55551
                56555
  50000
        9999
              0
1
                10
Press any key to continue . .
Number of (2x2) sub Matrices are 2
Number of (3x3) sub Matrices are 1
     5
        55
           ទាទាទាទ
             ຣຣຣຣ
                জ জে জে জ
9 10 10 10 10
  555555
     555
        SSS
Total Dominator Chromatic Number is
 2 + [(2 * number of (2x2) matrices) + (3 * number of (3x3) matrices)]
TOTAL DOMINATOR CHROMATIC NUMBER IS 9
```

REFERENCES

- 1. Harrary F, Graph theory, Addition-Wesley Reading Mass. 1969.
- Vijayalekshmi.A, Total dominator colorings in paths, International Journal of Mathematical Combinatorics, 2012; 2: 89-95

- Vijayalekshmi.A, Virgin Alangara Sheeba.J, Total dominator chromatic number of Paths, Cycles and Ladder graphs, International Journal of Con temporary Mathematical Sciences, Vol 13,2018,no. 5,199-204
- Vijayalekshmi.A, Total dominator colorings in cycles, International Journal of Mathematical Combinatorics, 2012; 4: 92-97
- 5. Jinnah M.I and Vijayalekshmi A, Total dominator colorings in graphs, Ph.D Thesis, University of Kerala. 2010
- 6. Terasa W.Haynes, Stephen T.Hedetniemi, Peter J.slater, Domination in Graphs, Marcel Dekker, New York, 1998
- 7. Terasa W.Haynes, Stephen T.Hedetniemi, Peter J.slater, Domination in Graphs Advanced Topics, Marceel Dekker, New York, 1998.